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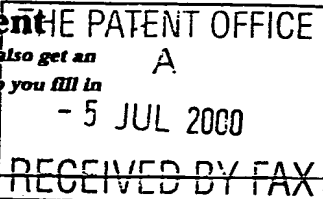
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P01/7700 0.00-0016475.6

Request for grant of a patent

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1. Your reference	2000P04864/GB/R76/MM/rr
2. Patent application number *(The Patent Office will fill in this part)*	**0016475.6** 0 5 JUL 2000

3. Full name, address and postcode of the or of each applicant *(underline all surnames)*

ROKE MANOR RESEARCH LIMITED
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UNITED KINGDOM

Patents ADP number *(if you know it)*

5615455006

If the applicant is a corporate body, give the country/state of its incorporation

UNITED KINGDOM

4. Title of the invention

DISTRIBUTED SCHEDULES FOR A ROUTING DEVICE

5. Name of your agent *(if you have one)*

DEREK ALLEN

"Address for service" in the United Kingdom to which all correspondence should be sent *(including the postcode)*

Siemens Shared Services Limited
Intellectual Property Department
Siemens House, Oldbury
Bracknell, Berkshire RG12 8FZ
United Kingdom

Patents ADP number *(if you know it)*

7396419002

6. If you are declaring priority from one or more earlier patent applications, give the country and the date of filing of the or of each of these earlier applications and *(if you know it)* the or each application number

Country	Priority application number *(if you know it)*	Date of filing *(day / month / year)*

7. If this application is divided or otherwise derived from an earlier UK application, give the number and the filing date of the earlier application

Number of earlier application	Date of filing *(day / month / year)*

8. Is a statement of inventorship and of right to grant of a patent required in support of this request? *(Answer 'Yes' if:*
a) *any applicant named in part 3 is not an inventor, or*
b) *there is an inventor who is not named as an applicant, or*
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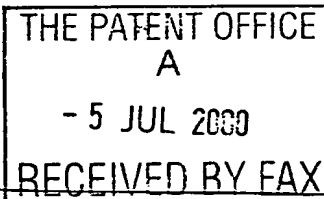Do not count copies of the same document

Continuation sheets of this form	0
Description	8
Claim (s)	0
Abstract	0
Drawing (s)	3

10. If you are also filing any of the following,
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Priority documents

Translations of priority documents

Statement of inventorship and right
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Request for preliminary examination
and search (Patents Form 9/77)

Request for substantive examination
(Patents Form 10/77)

Any other documents
(please specify)

11.

I/We request the grant of a patent on the basis of this application.

Signature _Magaret MACKETT_   Date   05.07.2000

12. Name and daytime telephone number of
person to contact in the United Kingdom

Margaret Mackett - 01344 396808

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Patents Form 1/77

DISTRIBUTED SCHEDULES FOR A ROUTING DEVICE

The present invention relates to a system for providing distributed schedules for a routing device. In particular, the invention relates to the

5 distribution of control between a management card and line interface cards (LICs) on a routing device. For the purposes of the following description the term routing device refers to any device which performs the function of a circuit switch or a router.

Data is transferred over the Internet by means of a plurality of

10 routing devices in accordance with a standard protocol known as Internet Protocol (IP). IP is a protocol based on the transfer of data in variable sized portions known as packets. All network traffic involves the transportation of packets of data. Routers are devices for accepting incoming packets; temporarily storing each packet; and then forwarding the

15 packets to another part of the network.

Traffic volume in the Internet is growing exponentially, almost doubling every 3 months, and the capacity of conventional IP routers is insufficient to meet this demand. There is thus an urgent requirement for products that can route IP traffic at extremely large aggregate bandwidths

20 in the order of several terabits per second. Such routing devices are termed "terabit routers".

Terabit routers require a scalable high capacity communications path between the point at which packets arrive at the router (the "ingress") and the point at which the packets leave the router (the "egress"). The terabit

25 routers discussed hereafter are intrinsically optical routing devices, since light is an effective and efficient medium for such high capacity communications.

The packets transferred in accordance with IP can (and do) vary in size. Within routers it has been found useful to pass data in fixed sized units. In routers then data packets are partitioned into small fixed sized units, known as cells. In discussion of routers, cells are often subdivided

5   into units of slices (64 bits) or octets (8 bits). A typical cell size is 64 octets, however cells of 128 octets are also known.

One suitable technique for implementing a scalable communications path is an optical backplane device, known as a cell based cross-bar. Data packets are partitioned into cells by a plurality of ingress line function

10   means for passage across the cross-bar.

The plurality of ingress line function means provide respective interfaces between incoming communications channels carrying incoming data and the optical backplane. Similarly a plurality of egress line function means provide respective interfaces between the optical backplane and

15   outgoing communications channels carrying outgoing data.

A general terabit router architecture bears some similarity to conventional router architecture. Packets of data arrive at input port(s) of ingress line function means and are routed as cells across the cross-bar to a predetermined egress line function means which reassembles the packets

20   and transmits them across its output port(s). Each ingress line function means maintains a separate packet queue for each egress line function means.

Line function means are implemented as line interface cards (LICs). Since one of the line functions regularly undertaken by line function means

25   is forwarding, LICs are also known as 'forwarders'. Further line functions include congestion control and maintenance of external interfaces, input ports and output ports.

In a conventional cell based cross-bar each ingress line function means is connected to one or more of the egress line function means. However, each ingress line function means is only capable of connecting to one egress line function means at any one time. Likewise, each egress line
5   function means is only capable of connecting to one ingress line function means at a time.

All ingress line function means transmit in parallel and independently across the cross-bar. Furthermore cell transmission is synchronised with a cell cycle, having a period of, for example, 108.8ns.

10   The ingress line function means simultaneously each transmit a new cell with each new cell cycle.

The pattern of transmissions from the ingress line function means across the cross-bar to the egress line function means changes at the end of every cell cycle.

15   The co-ordination of the transmission and reception of cells is performed by a cross-bar controller.

A cross-bar controller is provided for efficient allocation of the bandwidth across the cross-bar. It calculates the rates that each ingress line function means must transmit to each egress line function means. This is
20   the same as the rate at which data must be transmitted from each packet queue. The calculation makes use of real-time information, including traffic measurements and indications from the ingress line function means. The indications from the ingress line function means include monitoring the current rates, queue lengths and buffer full flags. The details of the
25   calculation are discussed more rigorously in the copending UK Patent Application Number 9907313.2 (docket number F21558/98P4863).

The cross-bar controller performs a further task; it serves to schedule the transfer of data efficiently across the cross-bar whilst maintaining the

calculated rates. At the end of each cell cycle, the cross-bar controller communicates with the ingress and egress line function means as follows. Firstly, the cross-bar controller calculates and transmits to each ingress line function means the identity of the next packet queue from which to

5   transmit. Secondly, the cross-bar controller calculates and transmits to each egress line function means the identity of the ingress from which it must receive.

The system described above does have a number of disadvantages however. The cross-bar controller is responsible for controlling the cell

10   cycle-by-cell cycle behaviour of each ingress and egress line function means. At the rates required by a terabit router, this amounts to demanding complex hardware to implement the crossbar controller, the ingress and the egress line function means. Furthermore the demand for higher capacity places stringent delay performance conditions upon the communication

15   channels between line function means and the cross-bar controller means.

When developing a system for particular traffic conditions, it is disadvantageous to have to replace inappropriate hardware.

It is therefore an object of the invention to obviate or at least mitigate the aforementioned problems.

20   In accordance with one aspect of the present invention, there is provided a routing device having a plurality of ingress line function means, a plurality of egress line function means, a backplane and a management card; each ingress line function means having: a schedule storing means, for storing a schedule of egress line function means addresses; a pointer

25   storing means, for storing a pointer to an address held in the schedule storing means; and a queue storing means, for storing a plurality of ingress queues of cells for transmission across the backplane, each one of the plurality of ingress queues corresponding uniquely to a predetermined one

of the egress line function means; wherein the management card communicates configuration primitives to each of the plurality of ingress line function means and to each of the plurality of egress line function means, the configuration primitives providing updated entries for the

5   schedule.

The ingress line function means may be line interface cards. Likewise the egress line function means may also be line interface cards.

As a result of the present invention, simple schedules for cell transmission and reception across the cross-bar are sent to each LIC by the

10   controller. The controller is responsible for formulating the schedules to ensure that traffic and quality of service requirements are met. The LICs merely have to obey the schedules, a simple task. The principle benefits are reduced complexity of the LICs, and flexibility. New services can be added to the router by software download without impact on the LIC

15   hardware.

For a better understanding of the present invention, reference will now be made, by way of example only, to the accompanying drawings in which:-

Figure 1 illustrates a terabit router architecture;

20   Figure 2 shows a cross-bar controller; and

Figure 3 shows a system for distributing schedules according to the present invention.

A conventional terabit router architecture is depicted in Figure 1. Packets of data arrive at ingress line function means (ingress forwarders)

25   via their input port(s) and are routed across the cross-bar to the correct egress line function means (egress forwarders) which transmits them across its output port(s). Each ingress line function means maintains a separate

packet queue for each egress line function means, for example q_{11}, q_{12} and q_{13}.

A conventional cell based cross-bar is shown in Figure 2. Here each ingress forwarder may be connected to one or more of the egress forwarders. However each ingress forwarder may only be connected to one egress forwarder at a time and each egress forwarder may only be connected to one ingress forwarder at a time.

The co-ordination of the transmission and reception of cells is performed by a cross-bar controller.

Each ingress forwarder communicates traffic measurements and notifications for the use of the cross-bar controller. The cross-bar controller allocates connections between ingress and egress forwarders and informs the respective forwarders accordingly for each cell cycle in turn.

Just as in Figure 2, the system in Figure 3 includes a backplane, a management card, which includes a cross-bar controller means, a plurality of ingress LICs and a plurality of egress LICs. The operation of the method according to the present invention will now be described using the system shown in Figure 3.

Each ingress LIC is associated with a respective schedule or timetable governing the transmission of cells by said ingress LIC. The schedule is in the form of a table whose entries are the identities of transmission queues, each queue corresponding to a respective egress LIC identification number. Each ingress LIC maintains a pointer into the schedule. At each cell transmit time, the LIC transmits a cell from the queue identified in the entry referenced by the pointer, and move the pointer to the next location. The schedule is circular in the sense that when moving the pointer from the last entry, its next position is the first entry.

Each egress LIC is also associated with a respective schedule governing the reception of cells by said egress LIC. The schedule is in the form of a table whose entries are the identities of ingress LICs from which to receive. An egress LIC maintains a pointer into the schedule. At each

5 cell transmit time, the LIC will receive the cell from the ingress identified in the entry referenced by the pointer, and move the pointer to the next location. Again, the schedule is circular in the sense that when moving the pointer from the last entry, its next position is the first entry.

The management card manages the contents of the ingress and egress

10 schedules. For each ingress LIC, the management card calculates the backplane rates required to each egress LIC. The backplane rates are calculated according to the current traffic load and required quality of service. Having calculated the rates, the management card calculates corresponding ingress and egress schedules needed to satisfy them. It then

15 updates the schedules using update messages known as configuration primitives. Only modifications to the schedules need be transmitted.

On receipt of the configuration primitives, the LICs update the schedules as requested.

This approach has the following advantages.

20 Firstly, complex rate calculation and schedule calculation algorithms can be implemented in software or on programmable devices on the management card, reducing development risk and cost.

Secondly, the hardware required to support the approach is very simple.

25 Thirdly, the communication channels between the LICs and the management card do not require stringent delay performance since the primitives are not directly synchronised to the cell transmissions across the backplane.

Fourthly, the communication channels between the LICs and the management card require less bandwidth than would be the case if the management card were to reconfigure the cross-bar every cell time.

Fifthly, the behaviour of the switch/router can be modified by changes to the software algorithms running on the management card. This has two consequences. Different applications can be supported by the same hardware. For example the same hardware could support both IP routing and circuit switching, including asynchronous transfer mode (ATM) and the related synchronous transfer mode (STM). Furthermore, optimisations and refinements to the algorithms can easily be implemented.

It will be readily understood that although the preceding discussion has been in terms of optical terabit routers, the apparatus of the present invention are capable of implementation in a wide variety of routing devices, including switches and routers, and that these routing devices can be either purely electronic, part electronic/part optical or optical in nature.

Further considerations when implementing the present invention, include the need for a degree of redundancy. To prevent a potentially catastrophic failure of a routing device, an further auxiliary management card having identical features to the management card can be provided to replace the management card if it were to fail. Since the operation of the routing device of the present invention is essentially stateless, the auxiliary management card can replace the management card without substantial interruption.
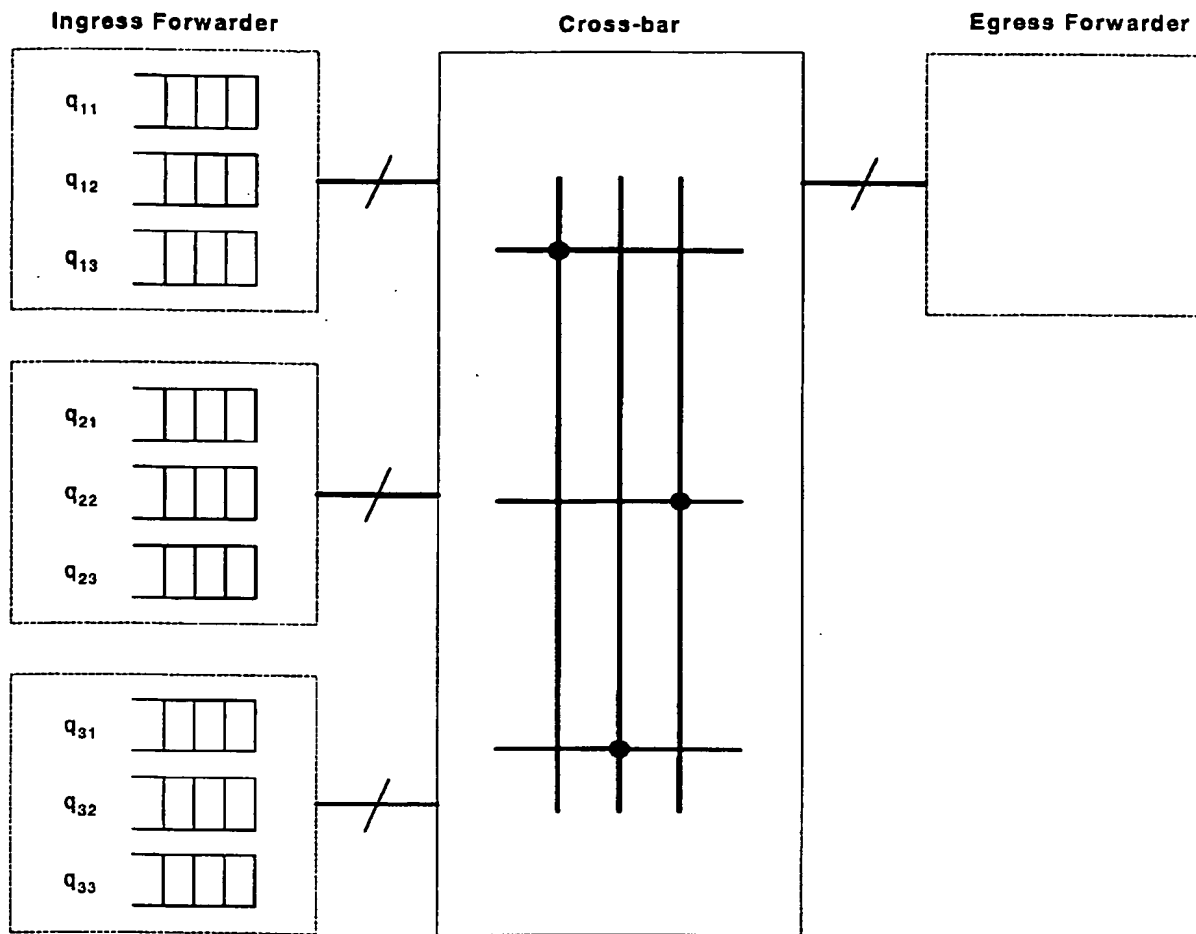
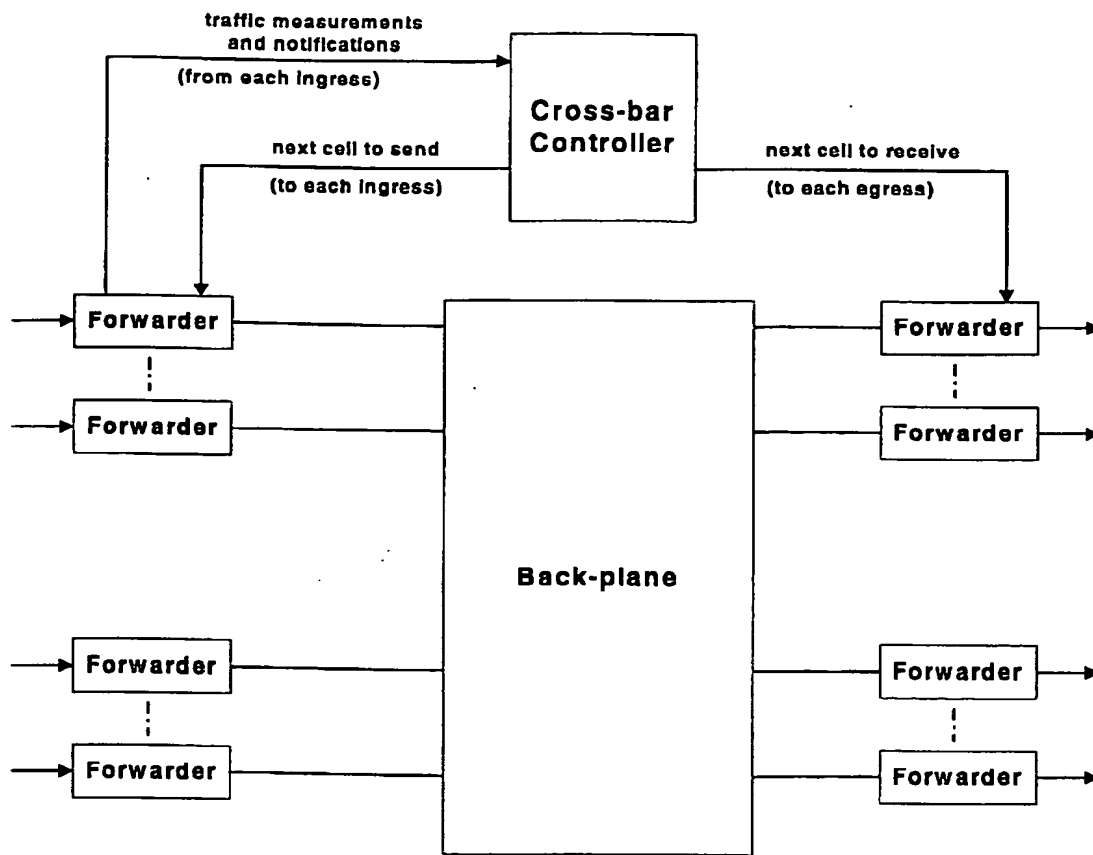**Figure 1: Terabit Router Architecture**
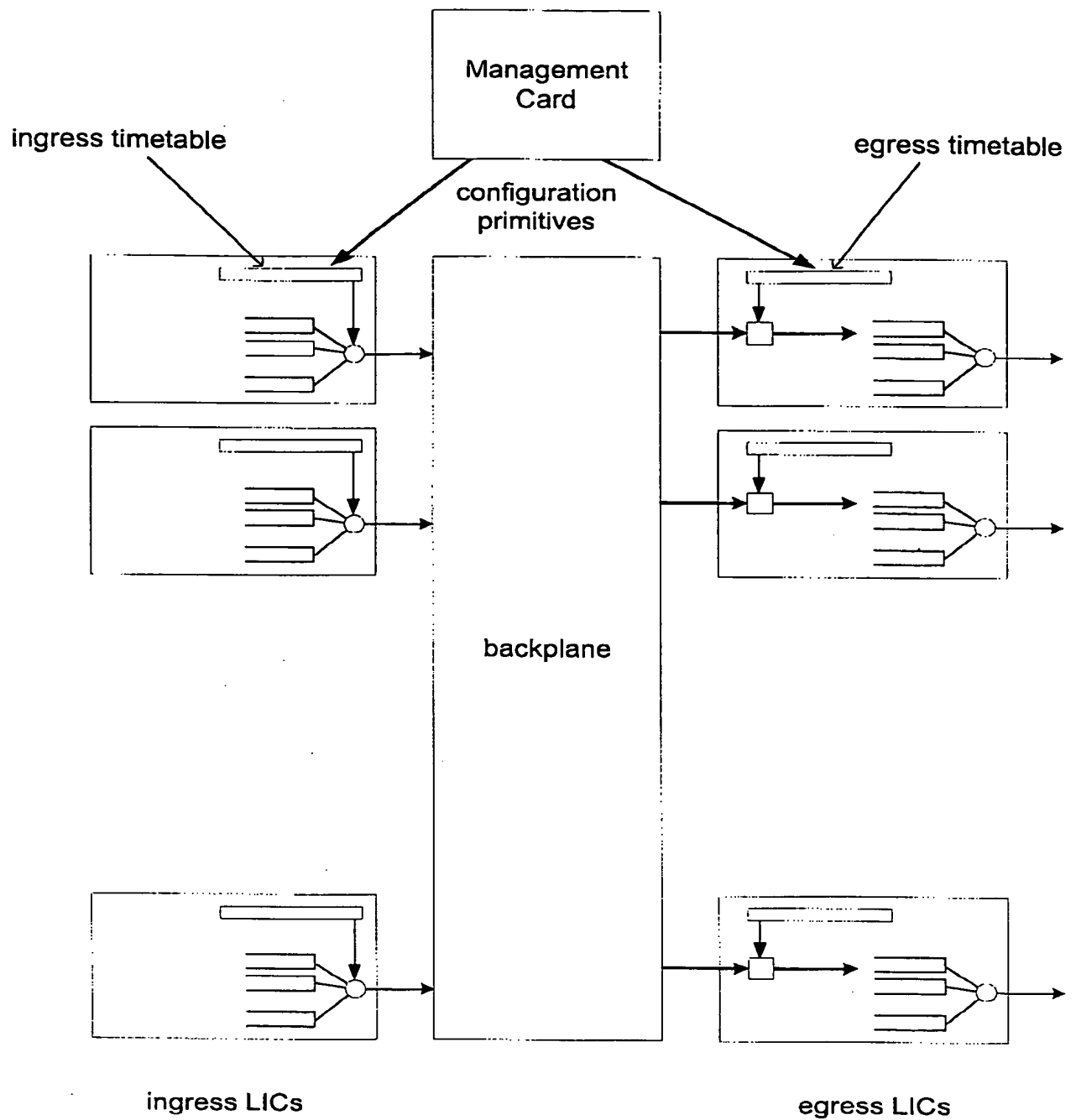
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Figure 2: Cross-bar Controller

Figure 3: Distributed Schedules

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